# Computable type of certain quotient spaces 

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## Computable topological spaces

Computable topological space is a second countable Hausdorff space ( $X, \mathcal{T}$ ) equipped with an enumeration of its base $\left(I_{i}\right)_{i \in \mathbb{N}}$ such that there exist r.e. sets $\mathcal{C}$ and $\mathcal{D}$ with the following properties:
i. $\mathcal{C} \subseteq\left\{(i, j) \mid I_{i} \subseteq I_{j}\right\}$;
ii. $\mathcal{D} \subseteq\left\{(i, j) \mid I_{i} \cap I_{j}=\emptyset\right\}$;
iii. $x \in I_{i} \cap I_{j} \Rightarrow \exists k\left(x \in I_{k} \wedge(k, i) \in \mathcal{C} \wedge(k, j) \in \mathcal{C}\right)$;
iv. $x \neq y \Rightarrow \exists i \exists j\left(x \in I_{i} \wedge y \in I_{j} \wedge(i, j) \in \mathcal{D}\right)$.

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- computable if $S$ is semicomputable and computably enumerable.


## Definition

Topological pair $(A, B)$ has computable type if, for every semicomputable copy $(S, T)$ of $(A, B)$ in a computable topological space, $S$ is computable.

## Computable type: examples


[Iljazović and Sušić, 2018, Čičković et al., 2019, Amir and Hoyrup, 2022, Horvat et al., 2020]

## Pairs and quotients



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## Definition

Let $B$ be a subset of a topological space $A$. The quotient space $A / B$ is the set $\{B\} \cup\{\{x\} \mid x \in A \backslash B\}$ equipped with the finest topology which makes the cannonical projection map continuous.

## The main goal



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## "Unbinding" a quotient space

## Theorem

Let $A$ be a topological space and let $B$ be a compact subset of $A$ such that $\operatorname{Int}_{A} B=\emptyset$. If $A / B$ has computable type, then $(A, B)$ also has computable type.

## "Unbinding" a quotient space

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Sketch of proof.

- Given a semicomputable set $B$ in a computable topological space, we can define a natural structure of computability on the corresponding quotient space
- The natural quotient map $q: X \rightarrow X / B$ preserves semicomputability
- The inverse image of $q$ preserves computable enumerability ( $\operatorname{Int}_{A} B=\emptyset$ is crucial in this step)

$$
\begin{gathered}
(A, B) \text { has } \\
\text { computable type }
\end{gathered} \stackrel{?}{\operatorname{Int}_{A} B=\emptyset} \quad \begin{gathered}
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## Counterexample I


$(A, B)$ has
computable type

## Counterexample I



## Counterexample I



Counterexample II

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## Counterexample II



## Local computability

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Suppose $T \subseteq S$. We say that $T$ is computably enumerable up to $S$ if there exists a r.e. set $\Omega \subseteq \mathbb{N}$ such that $\left\{i \in \mathbb{N} \mid I_{i} \cap T \neq \emptyset\right\} \subseteq \Omega \subseteq\left\{i \in \mathbb{N} \mid I_{i} \cap S \neq \emptyset\right\}$.

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We say that $S$ is computably enumerable at $x \in S$ if $x$ has a neighborhood in $S$ which is c.e. up to $S$.

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## Definition

Topological space $A$ has local computable type if a semicomputable set is c.e. at any point which has a neighborhood homeomorphic to $A$.

## Local computable type: examples





## Local computable type of $\mathbb{R}^{n} / K$

## Theorem

Let $K$ be a compact subset of $\mathbb{R}^{n}$ such that $\mathbb{R}^{n} \backslash K$ has finitely many connected components. Then $\mathbb{R}^{n} / K$ has local computable type.

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Sketch of proof.

- Suppose $S$ is semicomputable and $x_{0} \in S$ has a neighborhood homeomorphic to $\mathbb{R}^{n} / K$. WLOG we can assume $x_{0}$ is the image of the "collapsed" set $K$; otherwise, it has an Euclidean neighborhood.
- The image of any unbounded connected component of $\mathbb{R}^{n} \backslash K$ contains a neighborhood of $x_{0}$ which is c.e. up to $S$.
- The image of any connected component can be expressed as an image of an unbounded component under a similar map.

Example


Example


Example


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## Thank you!

