

Diagonalization Methods for Solving Definite Generalized Eigenvalue Problem

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OUTLINE

- GEP (DGEP, PGEP)

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- Global convergence of block algorithms

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For a definite pair (A, B) there exists a **nonsingular matrix** F such that

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Hence, probably the best choice for the kernel algorithm is some **element-wise diagonalization method**.

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The methods are connected: the FL method can be viewed as the HZ or CJ method with “fast scaled” transformations.

We have also derived their “equally promising” complex counterparts.

The Real and Complex FL Method

Starting with a **definite pair** (A, B) of Hermitian matrices, FL generates a sequence of “congruent” matrix pairs

$$(A, B) = (A^{(0)}, B^{(0)}), (A^{(1)}, B^{(1)}), \dots$$

by the rule

$$A^{(k+1)} = F_k^* A^{(k)} F_k, \quad B^{(k+1)} = F_k^* B^{(k)} F_k, \quad k \geq 0.$$

Here F_k is an **elementary plane matrix** defined by the **pivot pair** $(i(k), j(k))$

$$F_k = \begin{bmatrix} I & & & & \\ & 1 & & \alpha_k & \\ & & I & & \\ & \beta_k & & 1 & \\ & & & & I \end{bmatrix} \begin{matrix} i(k) \\ \\ j(k) \\ \\ \end{matrix}, \quad \alpha_k, \beta_k \in \mathbf{C},$$

Derivation of the Complex FL Method

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We have

$$A' = F^* A F, \quad B' = F^* B F \quad \left(\hat{A}' = \hat{F}^* \hat{A} \hat{F}, \quad \hat{B}' = \hat{F}^* \hat{B} \hat{F} \right)$$

and F is chosen to obtain $a'_{ij} = 0$ and $b'_{ij} = 0$.

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$$\begin{bmatrix} 1 & \bar{\beta} \\ \bar{\alpha} & 1 \end{bmatrix} \begin{bmatrix} a_{ii} & a_{ij} \\ \bar{a}_{ij} & a_{jj} \end{bmatrix} \begin{bmatrix} 1 & \alpha \\ \beta & 1 \end{bmatrix} = \begin{bmatrix} a'_{ii} & 0 \\ 0 & a'_{jj} \end{bmatrix}$$
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This leads us to solving a system of two nonlinear equations

$$\begin{aligned} e_1 &= a_{ii}\alpha + a_{jj}\bar{\beta} + \bar{a}_{ij}\alpha\bar{\beta} + a_{ij} = 0 \\ e_2 &= b_{ii}\alpha + b_{jj}\bar{\beta} + \bar{b}_{ij}\alpha\bar{\beta} + b_{ij} = 0. \end{aligned}$$

Solution if Matrices are Real and Symmetric

$$\mathfrak{S}_{ii} = a_{ii}b_{ij} - a_{ij}b_{ii} = \begin{vmatrix} a_{ii} & b_{ii} \\ a_{ij} & b_{ij} \end{vmatrix}$$

$$\mathfrak{S}_{jj} = a_{jj}b_{ij} - a_{ij}b_{jj} = \begin{vmatrix} a_{jj} & b_{jj} \\ a_{ij} & b_{ij} \end{vmatrix}$$

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$$\mathfrak{S} = \mathfrak{S}_{ij}^2 + 4\mathfrak{S}_{ii}\mathfrak{S}_{jj}$$

$$\nu = (\mathfrak{S}_{ij} + \operatorname{sgn}(\mathfrak{S}_{ij})\sqrt{\mathfrak{S}})/2$$

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If $\begin{bmatrix} a_{ii} & a_{ij} \\ \bar{a}_{ij} & a_{jj} \end{bmatrix}$ and $\begin{bmatrix} b_{ii} & b_{ij} \\ \bar{b}_{ij} & b_{jj} \end{bmatrix}$ are proportional, all \mathfrak{S}_{ii} , \mathfrak{S}_{jj} , \mathfrak{S}_{ij} , \mathfrak{S} and ν are zero and a special algorithm is required. [This is the real FL algorithm.](#)

The Complex FL Algorithm

$$\begin{aligned}\mathfrak{S}_{ii} &= a_{ii}b_{ij} - a_{ij}b_{ii} = \begin{vmatrix} a_{ii} & b_{ii} \\ a_{ij} & b_{ij} \end{vmatrix} \\ \mathfrak{S}_{jj} &= a_{jj}b_{ij} - a_{ij}b_{jj} = \begin{vmatrix} a_{jj} & b_{jj} \\ a_{ij} & b_{ij} \end{vmatrix} \\ \mathfrak{S}'_{ij} &= a_{ii}b_{jj} - a_{jj}b_{ii} = \begin{vmatrix} a_{ii} & b_{ii} \\ a_{jj} & b_{jj} \end{vmatrix} \\ \imath\mathfrak{S}''_{ij} &= a_{ij}\bar{b}_{ij} - \bar{a}_{ij}b_{ij} = \begin{vmatrix} a_{ij} & b_{ij} \\ \bar{a}_{ij} & \bar{b}_{ij} \end{vmatrix} = -2\imath \begin{vmatrix} \operatorname{Re}(a_{ij}) & \operatorname{Re}(b_{ij}) \\ \operatorname{Im}(a_{ij}) & \operatorname{Im}(b_{ij}) \end{vmatrix} \\ \mathfrak{S}_{ij} &= \mathfrak{S}'_{ij} + \imath\mathfrak{S}''_{ij} \\ \mathfrak{S} &= \mathfrak{S}_{ij}^2 + 4\bar{\mathfrak{S}}_{ii}\mathfrak{S}_{jj} = (\mathfrak{S}'_{ij})^2 - (\mathfrak{S}''_{ij})^2 + 2\imath\mathfrak{S}'_{ij}\mathfrak{S}''_{ij} + 4\bar{\mathfrak{S}}_{ii}\mathfrak{S}_{jj} \\ \nu &= (\mathfrak{S}_{ij} + \operatorname{sgn}(\mathfrak{S}'_{ij})\sqrt{\mathfrak{S}})/2, \\ \alpha &= \mathfrak{S}_j/\nu, \quad \beta = -\bar{\mathfrak{S}}_i/\nu \end{aligned}$$

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- **Difficult and challenging for making a good numerical code** (to many freedoms, all we have $\alpha A + \beta B \succ O$, when to stop iterations?)
- **Theoretical results are lacking** (all we have is quadratic asymptotic convergence result)

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The selection of pivot pairs $(i(k), j(k))$ defines pivot strategy.

At step k we denote: $A^{(k)} \mapsto A$, $A^{(k+1)} \mapsto A'$, $Z_k \mapsto Z$,

$$\hat{A} = \begin{bmatrix} a_{ii} & a_{ij} \\ \bar{a}_{ij} & a_{jj} \end{bmatrix}, \quad \hat{B} = \begin{bmatrix} 1 & b_{ij} \\ \bar{b}_{ij} & 1 \end{bmatrix}, \quad \hat{Z} = \begin{bmatrix} c & -s \\ \tilde{s} & \tilde{c} \end{bmatrix}.$$

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\hat{Z} is chosen to **diagonalize** \hat{A}' and **to make** \hat{B}' identity matrix I_2 .

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$$\hat{A} = \begin{bmatrix} a_{ii} & a_{ij} \\ \bar{a}_{ij} & a_{jj} \end{bmatrix}, \quad \hat{B} = \begin{bmatrix} 1 & b_{ij} \\ \bar{b}_{ij} & 1 \end{bmatrix}, \quad \hat{Z} = \begin{bmatrix} c & -s \\ \tilde{s} & \tilde{c} \end{bmatrix}.$$

\hat{A} , \hat{B} , \hat{Z} are **pivot submatrices** of A , B , Z .

Then $A' = Z^*AZ$, $B' = Z^*BZ$ implies $\hat{A}' = \hat{Z}^*\hat{A}\hat{Z}$, $\hat{B}' = \hat{Z}^*\hat{B}\hat{Z}$.

\hat{Z} is chosen to **diagonalize** \hat{A}' and **to make** \hat{B}' identity matrix I_2 .

\hat{Z} is sought in the form of a product of two Jacobi rotations and one or two diagonal matrices.

Real Algorithm: \hat{Z} is sought in the form:

$$(a) \begin{bmatrix} \frac{\sqrt{2}}{2} & -\frac{\sqrt{2}}{2} \\ \frac{\sqrt{2}}{2} & \frac{\sqrt{2}}{2} \end{bmatrix} \begin{bmatrix} \frac{1}{\sqrt{1+b_{ij}}} & 0 \\ 0 & \frac{1}{\sqrt{1-b_{ij}}} \end{bmatrix} \begin{bmatrix} \cos(\theta - \frac{\pi}{4}) & -\sin(\theta - \frac{\pi}{4}) \\ \sin(\theta - \frac{\pi}{4}) & \cos(\theta - \frac{\pi}{4}) \end{bmatrix}$$

$$(b) \begin{bmatrix} \frac{\sqrt{2}}{2} & \frac{\sqrt{2}}{2} \\ -\frac{\sqrt{2}}{2} & \frac{\sqrt{2}}{2} \end{bmatrix} \begin{bmatrix} \frac{1}{\sqrt{1-b_{ij}}} & 0 \\ 0 & \frac{1}{\sqrt{1+b_{ij}}} \end{bmatrix} \begin{bmatrix} \cos(\theta + \frac{\pi}{4}) & -\sin(\theta + \frac{\pi}{4}) \\ \sin(\theta + \frac{\pi}{4}) & \cos(\theta + \frac{\pi}{4}) \end{bmatrix}$$

$$\downarrow \\ \hat{B} \rightarrow \text{diag}$$

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 \text{(a)} \quad \begin{bmatrix} \frac{\sqrt{2}}{2} & -\frac{\sqrt{2}}{2} \\ \frac{\sqrt{2}}{2} & \frac{\sqrt{2}}{2} \end{bmatrix} \begin{bmatrix} \frac{1}{\sqrt{1+b_{ij}}} & 0 \\ 0 & \frac{1}{\sqrt{1-b_{ij}}} \end{bmatrix} \begin{bmatrix} \cos(\theta - \frac{\pi}{4}) & -\sin(\theta - \frac{\pi}{4}) \\ \sin(\theta - \frac{\pi}{4}) & \cos(\theta - \frac{\pi}{4}) \end{bmatrix} \\
 \text{(b)} \quad \begin{bmatrix} \frac{\sqrt{2}}{2} & \frac{\sqrt{2}}{2} \\ -\frac{\sqrt{2}}{2} & \frac{\sqrt{2}}{2} \end{bmatrix} \begin{bmatrix} \frac{1}{\sqrt{1-b_{ij}}} & 0 \\ 0 & \frac{1}{\sqrt{1+b_{ij}}} \end{bmatrix} \begin{bmatrix} \cos(\theta + \frac{\pi}{4}) & -\sin(\theta + \frac{\pi}{4}) \\ \sin(\theta + \frac{\pi}{4}) & \cos(\theta + \frac{\pi}{4}) \end{bmatrix} \\
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 \hat{B} \rightarrow \text{diag} \qquad \qquad \hat{B} \rightarrow I_2 \qquad \qquad \hat{A} \rightarrow \text{diag}
 \end{array}$$

Both approaches yield the same algorithm.

Essential Part of the Real Algorithm

$$\xi = \frac{b_{ij}}{\sqrt{1+b_{ij}} + \sqrt{1-b_{ij}}}, \quad \rho = \frac{1}{2}(\sqrt{1+b_{ij}} + \sqrt{1-b_{ij}}), \quad \xi^2 + \rho^2 = 1,$$

$$\tan(2\theta) = \frac{2a_{ij} - (a_{ii} + a_{jj}) b_{ij}}{\sqrt{1 - (b_{ij})^2} (a_{ii} - a_{jj})}, \quad -\frac{\pi}{4} \leq \theta \leq \frac{\pi}{4},$$

$$\cos \phi = \rho \cos \theta - \xi \sin \theta$$

$$\sin \phi = \rho \sin \theta + \xi \cos \theta$$

$$\cos \psi = \rho \cos \theta + \xi \sin \theta$$

$$\sin \psi = \rho \sin \theta - \xi \cos \theta$$

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$$\hat{Z} = \frac{1}{\sqrt{1 - b_{ij}^2}} \begin{bmatrix} \cos \phi & -\sin \phi \\ \cos \psi & \sin \psi \end{bmatrix}.$$

$$a'_{ii} = a_{ii} + \frac{1}{1 - b_{ij}^2} [(b_{ij}^2 - \sin^2 \phi) a_{ii} + 2 \cos \phi \sin \psi a_{ij} + \sin^2 \psi a_{jj}]$$

$$a'_{jj} = a_{jj} - \frac{1}{1 - b_{ij}^2} [(\sin^2 \psi - b_{ij}^2) a_{jj} + 2 \cos \psi \sin \phi a_{ij} + \sin^2 \phi a_{ii}]$$

Complex Algorithm: \hat{Z} is sought in the form:

$$\begin{array}{ccc}
 \hat{B} \rightarrow \text{diag} & & \hat{B} \rightarrow I_2 \\
 \uparrow & & \uparrow \\
 \hat{Z} = \begin{bmatrix} \frac{\sqrt{2}}{2} & -\frac{\sqrt{2}}{2} e^{i \arg(b_{ij})} \\ \frac{\sqrt{2}}{2} e^{-i \arg(b_{ij})} & \frac{\sqrt{2}}{2} \end{bmatrix} \cdot \begin{bmatrix} \frac{1}{\sqrt{1+|b_{ij}|}} & 0 \\ 0 & \frac{1}{\sqrt{1-|b_{ij}|}} \end{bmatrix} \\
 \cdot \begin{bmatrix} \cos(\theta - \frac{\pi}{4}) & -e^{i\alpha} \sin(\theta - \frac{\pi}{4}) \\ e^{-i\alpha} \sin(\theta - \frac{\pi}{4}) & \cos(\theta - \frac{\pi}{4}) \end{bmatrix} \cdot \begin{bmatrix} e^{i\omega_j} & 0 \\ 0 & e^{i\omega_j} \end{bmatrix} \\
 \downarrow & & \downarrow \\
 \hat{A} \rightarrow \text{diag} & & \text{diag}(\hat{Z}) \succ 0
 \end{array}$$

Essential Part of the Complex Algorithm

Let

$$b = |b_{ij}|, \quad t = \sqrt{1 - b^2}, \quad e = a_{jj} - a_{ii}, \quad \epsilon = \begin{cases} 1, & e \geq 0 \\ -1, & e < 0 \end{cases},$$

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$$u + \imath v = e^{-\imath \arg(b_{ij})} a_{ij}, \quad \tan \gamma = 2 \frac{v}{e}, \quad -\frac{\pi}{2} < \gamma \leq \frac{\pi}{2}$$

$$\tan 2\theta = \epsilon \frac{2u - (a_{ii} + a_{jj})b}{t\sqrt{e^2 + 4v^2}}, \quad -\frac{\pi}{4} < \theta \leq \frac{\pi}{4}$$

$$2 \cos^2 \phi = 1 + b \sin 2\theta + t \cos 2\theta \cos \gamma, \quad 0 \leq \phi \leq \frac{\pi}{2}$$

$$2 \cos^2 \psi = 1 - b \sin 2\theta + t \cos 2\theta \cos \gamma, \quad 0 \leq \psi \leq \frac{\pi}{2}$$

$$e^{\imath\alpha} \sin \phi = \frac{e^{\imath \arg(b_{ij})}}{2 \cos \psi} [\sin 2\theta - b - \imath t \cos 2\theta \sin \gamma]$$

$$e^{-\imath\beta} \sin \psi = \frac{e^{-\imath \arg(b_{ij})}}{2 \cos \phi} [\sin 2\theta + b + \imath t \cos 2\theta \sin \gamma].$$

Then

$$\hat{Z} = \frac{1}{\sqrt{1 - b^2}} \begin{bmatrix} \cos \phi & e^{\imath\alpha} \sin \phi \\ -e^{-\imath\beta} \sin \psi & \cos \psi \end{bmatrix}$$

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- The real CJ algorithm is obtained by simplifying the complex one

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Consider the **Cholesky factorization** of \hat{B} : $\hat{B} = \hat{L}\hat{L}^*$,

$$\begin{bmatrix} 1 & b_{ij} \\ \bar{b}_{ij} & 1 \end{bmatrix} = \hat{B} = \hat{L}\hat{L}^* = \begin{bmatrix} 1 & 0 \\ \bar{a} & \bar{c} \end{bmatrix} \begin{bmatrix} 1 & a \\ 0 & c \end{bmatrix} = \begin{bmatrix} 1 & a \\ \bar{a} & |a|^2 + |c|^2 \end{bmatrix}$$

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Let $\hat{F}_1 = \hat{L}^{-*}$. Then $\hat{F}_1^* \hat{B} \hat{F}_1 = I_2$ and

$$\hat{F}_1^* \hat{A} \hat{F}_1 = \begin{bmatrix} a_{ii} & (a_{ij} - b_{ij}a_{ii})/\tau \\ (\bar{a}_{ij} - \bar{b}_{ij}a_{ii})/\tau & a_{jj} - \frac{a_{ij}\bar{b}_{ij} + \bar{a}_{ij}b_{ij} - (a_{ii} + a_{jj})|b_{ij}|^2}{1 - |b_{ij}|^2} \end{bmatrix}.$$

Derivation of Complex LL^*J Algorithm, $B \succ O$

The final \hat{F} is obtained as product $\hat{F} = \hat{F}_1 \hat{R}_1$ where

\hat{R}_1 is the complex Jacobi rotation which diagonalizes $\hat{F}_1^* \hat{A} \hat{F}_1$.

Let us assume that the $(1, 2)$ -element of \hat{R}_1 is $-e^{2\epsilon_1} \sin \vartheta_1$. Then the angles ϑ_1 and ϵ_1 are determined by the formulas

$$\begin{aligned} \epsilon_1 &= \arg(a_{ij} - b_{ij}a_{ii}), \\ \tan(2\vartheta_1) &= \frac{2|a_{ij} - a_{ii}b_{ij}|\sqrt{1 - |b_{ij}|^2}}{a_{ii} - a_{jj} + a_{ij}\bar{b}_{ij} + \bar{a}_{ij}b_{ij} - 2a_{ii}|b_{ij}|^2}, \quad -\frac{\pi}{4} \leq \vartheta_1 \leq \frac{\pi}{4}. \end{aligned}$$

Derivation of Complex LL^*J Algorithm, $B \succ O$

The transformation formulas for the diagonal elements of A read

$$a'_{ii} = a_{ii} + \tan \vartheta_1 \cdot \frac{|a_{ij} - a_{ii}b_{ij}|}{\sqrt{1 - b_{ij}^2}}$$
$$a'_{jj} = a_{jj} - \frac{a_{ij}\bar{b}_{ij} + \bar{a}_{ij}b_{ij} - (a_{ii} + a_{jj})|b_{ij}|^2}{1 - |b_{ij}|^2} - \tan \vartheta_1 \cdot \frac{|a_{ij} - a_{ii}b_{ij}|}{\sqrt{1 - b_{ij}^2}}.$$

In the case $a_{ii} = a_{jj}$, $a_{ij} = a_{ii}b_{ij}$, $\tan(2\vartheta_1)$ has the form $0/0$.

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In the case $a_{ii} = a_{jj}$, $a_{ij} = a_{ii}b_{ij}$, $\tan(2\vartheta_1)$ has the form $0/0$.

Then we choose $\vartheta_1 = 0$, so that $a'_{ii} = a_{ii}$ and $a'_{jj} = a_{jj}$.

Derivation of Complex LL^*J Algorithm, $B \succ O$

Let $c_{\vartheta_1} = \cos \vartheta_1$, $s_{\vartheta_1}^{\pm} = e^{\pm i\epsilon_1} \sin \vartheta_1$. Then

$$\begin{aligned}\hat{F} &= \frac{1}{\sqrt{1 - |b_{ij}|^2}} \begin{bmatrix} \sqrt{1 - |b_{ij}|^2} & -b_{ij} \\ 0 & 1 \end{bmatrix} \begin{bmatrix} c_{\vartheta_1} & -s_{\vartheta_1}^+ \\ s_{\vartheta_1}^- & c_{\vartheta_1} \end{bmatrix} \\ &= \frac{1}{\sqrt{1 - |b_{ij}|^2}} \begin{bmatrix} c_{\tilde{\vartheta}_1} & -s_{\tilde{\vartheta}_1} \\ s_{\vartheta_1}^- & c_{\vartheta_1} \end{bmatrix} \quad \begin{aligned} c_{\tilde{\vartheta}_1} &= c_{\vartheta_1} \sqrt{1 - |b_{ij}|^2} - s_{\vartheta_1}^- b_{ij} \\ s_{\tilde{\vartheta}_1} &= c_{\vartheta_1} b_{ij} + s_{\vartheta_1}^+ \sqrt{1 - |b_{ij}|^2} \end{aligned} \\ &= \begin{bmatrix} c1 & -s1 \\ s2 & c2 \end{bmatrix},\end{aligned}$$

$$\begin{aligned}c1 &= c_{\vartheta_1} - s_{\vartheta_1}^- b_{ij} / \sqrt{1 - |b_{ij}|^2}, & c2 &= c_{\vartheta_1} / \sqrt{1 - |b_{ij}|^2}, \\ s1 &= c_{\vartheta_1} b_{ij} / \sqrt{1 - |b_{ij}|^2} + s_{\vartheta_1}^+, & s2 &= s_{\vartheta_1}^- / \sqrt{1 - |b_{ij}|^2}.\end{aligned}$$

Derivation of Complex RR^*J Algorithm, $B \succ O$

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$$\begin{bmatrix} 1 & b_{ij} \\ \bar{b}_{ij} & 1 \end{bmatrix} = \hat{B} = \hat{R}\hat{R}^* = \begin{bmatrix} c & a \\ 0 & 1 \end{bmatrix} \begin{bmatrix} \bar{c} & 0 \\ \bar{a} & 1 \end{bmatrix} = \begin{bmatrix} |a|^2 + |c|^2 & a \\ \bar{a} & 1 \end{bmatrix}.$$

Assuming positive c , one obtains $a = b_{ij}$, $c = \sqrt{1 - |b_{ij}|^2} = \tau$. Hence

$$\hat{R} = \begin{bmatrix} \tau & b_{ij} \\ 0 & 1 \end{bmatrix}, \quad \hat{R}^{-1} = \frac{1}{\tau} \begin{bmatrix} 1 & -b_{ij} \\ 0 & \tau \end{bmatrix}, \quad \hat{R}^{-*} = \frac{1}{\tau} \begin{bmatrix} 1 & 0 \\ -\bar{b}_{ij} & \tau \end{bmatrix}.$$

If we write $\hat{F}_2 = \hat{R}^{-*}$, then $\hat{F}_2^* \hat{B} \hat{F}_2 = \hat{R}^{-1} \hat{B} \hat{R}^{-*} = I_2$ and we have

The Algorithm RR^*J

$$\hat{F}_2^* \hat{A} \hat{F}_2 = \begin{bmatrix} a_{ii} - \frac{a_{ij}\bar{b}_{ij} + \bar{a}_{ij}b_{ij} - (a_{ii} + a_{jj})|b_{ij}|^2}{\tau^2} & (a_{ij} - a_{jj}b_{ij})/\tau \\ (\bar{a}_{ij} - a_{jj}\bar{b}_{ij})/\tau & a_{jj} \end{bmatrix}.$$

- The final transformation is $\hat{F} = \hat{F}_2 \hat{R}_2$,
- \hat{R}_2 is the Jacobi rotation which annihilates $(1, 2)$ -element of $\hat{F}_2^* \hat{A} \hat{F}_2$
- Let $(1, 2)$ -element of \hat{R}_2 be $-e^{i\epsilon_2} \sin \vartheta_2$

Then the parameters ϵ_2 and ϑ_2 are determined by the formulas

$$\begin{aligned} \epsilon_2 &= \arg(a_{ij} - b_{ij}a_{jj}), \\ \tan(2\vartheta_2) &= \frac{2|a_{ij} - a_{jj}b_{ij}|\sqrt{1 - |b_{ij}|^2}}{a_{ii} - a_{jj} - (a_{ij}\bar{b}_{ij} + \bar{a}_{ij}b_{ij}) + 2a_{jj}|b_{ij}|^2}, \quad -\frac{\pi}{4} \leq \vartheta_2 \leq \frac{\pi}{4}. \end{aligned}$$

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$$a'_{jj} = a_{jj} - \tan \vartheta_2 \cdot \frac{|a_{ij} - a_{jj}b_{ij}|}{\sqrt{1 - b_{ij}^2}}.$$

If $a_{ii} = a_{jj}$, $a_{ij} = a_{jj}b_{ij}$, ϑ_2 is not well defined and we choose $\vartheta_2 = 0$.

In that case a'_{ii} and a'_{jj} reduce to a_{ii} and a_{jj} , respectively.

The Algorithm RR^*J

Let $c_{\vartheta_2} = \cos \vartheta_2$, $s_{\vartheta_2}^{\pm} = e^{\pm i\epsilon_2} \sin \vartheta_2$. Then

$$\begin{aligned}\hat{F} &= \frac{1}{\sqrt{1 - |b_{ij}|^2}} \begin{bmatrix} 1 & 0 \\ -\bar{b}_{ij} & \sqrt{1 - |b_{ij}|^2} \end{bmatrix} \begin{bmatrix} c_{\vartheta_2} & -s_{\vartheta_2}^+ \\ s_{\vartheta_2}^- & c_{\vartheta_2} \end{bmatrix} \\ &= \frac{1}{\sqrt{1 - |b_{ij}|^2}} \begin{bmatrix} c_{\vartheta_2} & -s_{\vartheta_2}^+ \\ s_{\tilde{\vartheta}_2} & c_{\tilde{\vartheta}_2} \end{bmatrix}, \quad \begin{aligned} c_{\tilde{\vartheta}_2} &= c_{\vartheta_2} \sqrt{1 - |b_{ij}|^2} + s_{\vartheta_2}^+ \bar{b}_{ij} \\ s_{\tilde{\vartheta}_2} &= s_{\vartheta_2}^- \sqrt{1 - |b_{ij}|^2} - c_{\vartheta_2} \bar{b}_{ij} \end{aligned} \\ &= \begin{bmatrix} c1 & -s1 \\ s2 & c2 \end{bmatrix}, \quad \text{It is easy to check that } c_{\tilde{\vartheta}}^2 + s_{\tilde{\vartheta}}^2 = 1.\end{aligned}$$

$$\begin{aligned}c1 &= c_{\vartheta_2} / \sqrt{1 - b_{ij}^2}, & c2 &= c_{\vartheta_2} + s_{\vartheta_2}^+ \bar{b}_{ij} / \sqrt{1 - b_{ij}^2}, \\ s1 &= s_{\vartheta_2}^+ / \sqrt{1 - b_{ij}^2}, & s2 &= s_{\vartheta_2}^- - c_{\vartheta_2} \bar{b}_{ij} / \sqrt{1 - b_{ij}^2}.\end{aligned}$$

We can postmultiply \hat{F} by $\text{diag}(1, \bar{c}_{\tilde{\vartheta}_2}/|c_{\tilde{\vartheta}_2}|)$ provided that $c_{\tilde{\vartheta}_2} \neq 0$. This ensures that (the updated) \hat{F} has nonnegative diagonal elements.

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Our numerical tests show that **neither LL^*J nor RR^*J** is indicated as a **high relative accurate** algorithm on pairs of **well-behaved positive definite matrices**.

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Only the above definition warrants the **high relative accuracy of the algorithm** and it is in complete agreement with the behavior of the real *CJ* method.

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- **It requires B to be positive definite** (it solves PGEP)

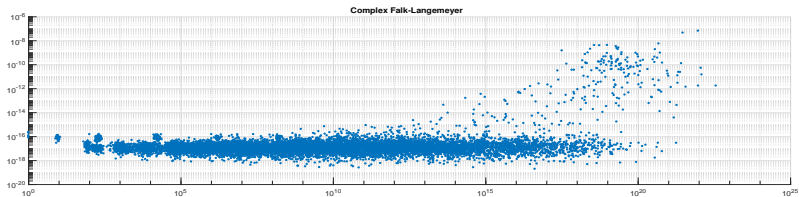
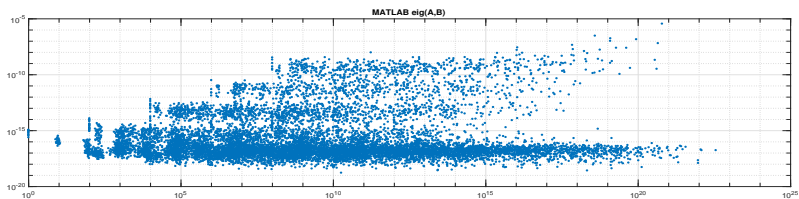
$$A_S = [\text{diag}(A)]^{-1/2} A [\text{diag}(A)]^{-1/2}, \quad B_S = [\text{diag}(B)]^{-1/2} B [\text{diag}(B)]^{-1/2}$$

$$\varrho_{(A,B)} = \max_{1 \leq i \leq n} \frac{|\tilde{\lambda}_i - \lambda_i|}{\lambda_i} / \sqrt{\kappa_2^2(A_S) + \kappa_2^2(B_S)}$$






$$\chi_{(A,B)} = \sqrt{\kappa_2^2(A^{(0)}) + \kappa_2^2(B^{(0)})}$$

$$\mathcal{E} = \{(\chi_{(A,B)}, \varrho_{(A,B)}) : (A, B) \in \mathcal{T}\}.$$

Relative errors: *CFL* vs. MATLAB eig(A,B)



Some References and Thank You for Your Patience!

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-  V. Novaković, S. Singer, S. Singer, *Blocking and Parallelization of the Hari-Zimmermann Variant of the Falk-Langemeyer Algorithm for the Generalized SVD*, Parallel Comput., 49 (2015) 136-152.