# On Jacobi Methods for the Positive Definite Generalized Eigenvalue Problem 

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- Derivation of the algorithms


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- Stability and relative accuracy

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- Block algorithms
- Global convergence of block algorithms
- We have restricted our attention to element-wise, two-sided Jacobi-type methods for PGEP since they can be used standalone or as kernel algorithms for the block methods.

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For such a pair $(A, B)$ there exists a nonsingular matrix $F$ such that
$F^{\top} A F=\Lambda_{A}=\operatorname{diag}\left(\alpha_{1}, \ldots, \alpha_{n}\right), \quad F^{\top} B F=\Lambda_{B}=\operatorname{diag}\left(\beta_{1}, \ldots, \beta_{n}\right) \succ O$,

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The eigenpairs of $(A, B)$ are: $\quad\left(\alpha_{i} / \beta_{i}, F e_{i}\right), \quad 1 \leq i \leq n ;$ where $I_{n}=\left[e_{1}, \ldots, e_{n}\right]$.

## How to solve PGEP?

One can reduce PGEP to the standard EP for one symmetric matrix

$$
(A, B) \mapsto\left(L^{-1} A L^{-T}, I\right), \quad B=L L^{T} .
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One can try to maximize the smallest eigenvalue of $B$ by rotating the pair

$$
(A, B) \mapsto\left(A_{\varphi}, B_{\varphi}\right)=(A \cos \varphi+B \sin \varphi,-A \sin \varphi+B \cos \varphi),
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or derive a method which works with the initial pair $(A, B)$.

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We follow the second path.

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- Falk-Langemeyer method (shorter: FL method)
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- Hari-Zimmermann variant of the FL method (shorter: HZ method) (Hari Ph.D. 1984)


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The two methods are connected: the FL method can be viewed as the HZ method with "fast scaled" transformations.

So, the FL method seems to be somewhat faster and the HZ method seems to be more robust.

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Since the derivation of the HZ method has not yet been published, we shall devote few slides to its derivation.

## Derivation of the HZ Method

Preliminary transformation: $\quad A^{(0)}=D_{0} A D_{0}, \quad B^{(0)}=D_{0} B D_{0}$

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This property of $B^{(0)}$ is maintained during the iteration process:

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Z_{k}=\left[\begin{array}{ccccc}
1 & & & & \\
& c_{k} & & -s_{k} & \\
& \tilde{s}_{k} & & \tilde{c}_{k} & \\
& & &
\end{array}\right] \begin{aligned}
& i(k) \\
& j(k)
\end{aligned}, \quad c_{k}^{2}+s_{k}^{2}=\tilde{c}_{k}^{2}+\tilde{s}_{k}^{2}=1 / \sqrt{1-b_{i(k) j(k)}^{2}},
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$$

The selection of pivot pairs $(i(k), j(k))$ defines pivot strategy.

## Derivation of the HZ Method

At step $k$ we denote: $\quad A^{(k)} \mapsto A, \quad A^{(k+1)} \mapsto A^{\prime}, \quad Z_{k} \mapsto Z$,

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\hat{A}=\left[\begin{array}{cc}
a_{i i} & a_{i j} \\
a_{i j} & a_{j j}
\end{array}\right], \quad \hat{B}=\left[\begin{array}{cc}
1 & b_{i j} \\
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\end{array}\right], \quad \hat{Z}=\left[\begin{array}{cc}
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$\hat{Z}$ is chosen to diagonalize $\hat{A}^{\prime}$ and to make $\hat{B}^{\prime}$ identity matrix $I_{2}$.
$\hat{Z}$ is sought in the form of a product of two Jacobi rotations and one diagonal matrix. We have two possibilities:

## $\hat{Z}$ is sought in the form:

(a) $\left[\begin{array}{cc}\frac{\sqrt{2}}{2} & -\frac{\sqrt{2}}{2} \\ \frac{\sqrt{2}}{2} & \frac{\sqrt{2}}{2}\end{array}\right]\left[\begin{array}{cc}\frac{1}{\sqrt{1+b_{j}}} & 0 \\ 0 & \frac{1}{\sqrt{1-b_{j}}}\end{array}\right]\left[\begin{array}{cc}\cos \left(\theta-\frac{\pi}{4}\right) & -\sin \left(\theta-\frac{\pi}{4}\right) \\ \sin \left(\theta-\frac{\pi}{4}\right) & \cos \left(\theta-\frac{\pi}{4}\right)\end{array}\right]$
(b) $\left[\begin{array}{cc}\frac{\sqrt{2}}{2} & \frac{\sqrt{2}}{2} \\ -\frac{\sqrt{2}}{2} & \frac{\sqrt{2}}{2}\end{array}\right]\left[\begin{array}{cc}\frac{1}{\sqrt{1-b_{i}}} & 0 \\ 0 & \frac{1}{\sqrt{1+b_{j}}}\end{array}\right]\left[\begin{array}{ll}\cos \left(\theta+\frac{\pi}{4}\right) & -\sin \left(\theta+\frac{\pi}{4}\right) \\ \sin \left(\theta+\frac{\pi}{4}\right) & \cos \left(\theta+\frac{\pi}{4}\right)\end{array}\right]$
$\hat{B} \rightarrow$ diag
$\hat{B} \rightarrow I_{2}$
$\hat{A} \rightarrow \operatorname{diag}$

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$\rightarrow$ diag

$\hat{B} \rightarrow$ diag
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The both approaches yield the same algorithm.

## Essential Part of the Algorithm

$$
\xi=\frac{b_{i j}}{\sqrt{1+b_{i j}}+\sqrt{1-b_{i j}}}, \quad \rho=\frac{1}{2}\left(\sqrt{1+b_{i j}}+\sqrt{1-b_{i j}}\right), \quad \xi^{2}+\rho^{2}=1,
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\begin{gathered}
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\tan (2 \theta)=\frac{2 a_{i j}-\left(a_{i i}+a_{j j}\right) b_{i j}}{\sqrt{1-\left(b_{i j}\right)^{2}}\left(a_{i i}-a_{j j}\right)}, \quad-\frac{\pi}{4} \leq \theta \leq \frac{\pi}{4},
\end{gathered}
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& \quad \tan (2 \theta)=\frac{2 a_{i j}-\left(a_{i i}+a_{j j}\right) b_{i j}}{\sqrt{1-\left(b_{i j}\right)^{2}}\left(a_{i i}-a_{j j}\right)}, \quad-\frac{\pi}{4} \leq \theta \leq \frac{\pi}{4} \\
& \cos \phi=\rho \cos \theta-\xi \sin \theta \\
& \sin \phi=\rho \sin \theta+\xi \cos \theta \\
& \cos \psi=\rho \cos \theta+\xi \sin \theta \\
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\cos \phi= & \rho \cos \theta-\xi \sin \theta \\
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\end{aligned} \quad \hat{Z}=\frac{1}{\sqrt{1-b_{i j}^{2}}}\left[\begin{array}{cc}
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& \hat{Z}=\frac{1}{\sqrt{1-b_{i j}^{2}}}\left[\begin{array}{cc}
\cos \phi & -\sin \phi \\
\cos \psi & \sin \psi
\end{array}\right] . \\
& a_{i i}^{\prime}=a_{i i}+\frac{1}{1-b_{i j}^{2}}\left[\left(b_{i j}^{2}-\sin ^{2} \phi\right) a_{i i}+2 \cos \phi \sin \psi a_{i j}+\sin ^{2} \psi a_{j j}\right] \\
& a_{j j}^{\prime}=a_{j j}-\frac{1}{1-b_{i j}^{2}}\left[\left(\sin ^{2} \psi-b_{i j}^{2}\right) a_{j j}+2 \cos \psi \sin \phi a_{i j}+\sin ^{2} \phi a_{i i}\right]
\end{aligned}
$$

## There are more formulas!

$$
2 \rho \xi=b_{i j}, \quad|\xi| \leq \sqrt{2} / 2 \leq \rho \leq 1 .
$$

$$
\begin{aligned}
\cos \phi \sin \psi & =\cos \theta \sin \theta-\rho \xi=0.5\left(\sin 2 \theta-b_{i j}\right) \\
\cos \psi \sin \phi & =\cos \theta \sin \theta+\rho \xi=0.5\left(\sin 2 \theta+b_{i j}\right) \\
\cos \phi \cos \psi & =\rho^{2} \cos ^{2} \theta-\xi^{2} \sin ^{2} \theta \\
\sin \phi \sin \psi & =\rho^{2} \sin ^{2} \theta-\xi^{2} \cos ^{2} \theta
\end{aligned}
$$

$$
\begin{aligned}
\min \{\cos \phi, \cos \psi\} & \geq \rho \cos \theta-\frac{\left|b_{i j}\right|}{2 \rho}|\sin \theta| \geq\left(\rho-\frac{\left|b_{i j}\right|}{2 \rho}\right) \cos \theta>0 \\
\max \{\cos \phi, \cos \psi\} & =\rho \cos \theta+|\xi \sin \theta| \geq \cos (\theta) \geq \frac{\sqrt{2}}{2}
\end{aligned}
$$

## There are more formulas!

Let

$$
\sin \gamma=b_{i j}, \quad \cos \gamma=\sqrt{1-b_{i j}^{2}}
$$

Then we have
$\frac{1}{\cos \gamma}\left[\begin{array}{ll}a_{i i} & a_{i j} \\ a_{i j} & a_{j j}\end{array}\right]\left[\begin{array}{cc}\cos \phi & -\sin \phi \\ \sin \psi & \cos \psi\end{array}\right]=\left[\begin{array}{cc}\cos \psi & -\sin \psi \\ \sin \phi & \cos \phi\end{array}\right]\left[\begin{array}{ll}a_{i i}^{\prime} & \\ & a_{j j}^{\prime}\end{array}\right]$,
$\frac{1}{\cos \gamma}\left[\begin{array}{cc}1 & b_{i j} \\ b_{i j} & 1\end{array}\right]\left[\begin{array}{cc}\cos \phi & -\sin \phi \\ \sin \psi & \cos \psi\end{array}\right]=\left[\begin{array}{cc}\cos \psi & -\sin \psi \\ \sin \phi & \cos \phi\end{array}\right]$,

$$
\begin{aligned}
\cos \gamma & =\frac{\cos \phi}{\cos \psi}+b_{i j} \tan \psi=\frac{\cos \psi}{\cos \phi}-b_{i j} \tan \phi \\
2 \cos (\phi+\psi) a_{i j} & =a_{i i} \sin (2 \phi)-a_{j j} \sin (2 \psi)
\end{aligned}
$$

## There are more formulas!

$$
\begin{aligned}
a_{i i}^{\prime} & =\frac{1}{\cos \gamma}\left(a_{i i} \frac{\cos \phi}{\cos \psi}+a_{i j} \tan \psi\right)=\frac{a_{i i}+a_{i j} \frac{\sin \psi}{\cos \phi}}{1+b_{i j} \frac{\sin \psi}{\cos \phi}} \\
a_{j j}^{\prime} & =\frac{1}{\cos \gamma}\left(a_{j j} \frac{\cos \psi}{\cos \phi}-a_{i j} \tan \phi\right)=\frac{a_{j j}-a_{i j} \frac{\sin \phi}{\cos \psi}}{1-b_{i j} \frac{\sin \phi}{\cos \psi}} .
\end{aligned}
$$

We also have

$$
\begin{aligned}
\phi+\psi & =2 \theta \\
\phi-\psi & =\gamma
\end{aligned}, \quad \text { hence } \quad \begin{aligned}
& \phi=\theta+\gamma / 2 \\
& \psi
\end{aligned}=\theta-\gamma / 2 .
$$

All these relations are used in the global convergence proof and in the proof of high relative accuracy of the method.

## Algorithm HZ

select the pivot pair $(i, j)$
if $a_{i j} \neq 0$ or $b_{i j} \neq 0$ then

$$
\begin{aligned}
& \rho=0.5\left(\sqrt{1+b_{i j}}+\sqrt{1-b_{i j}}\right) ; \quad \xi=b_{i j} /(2 \rho) ; \\
& \tau=\sqrt{\left(1+b_{i j}\right)\left(1-b_{i j}\right) ; \quad t 2=2 a_{i j}-\left(a_{i i}+a_{i j}\right) b_{i j} ;} \\
& \text { if } t 2=0 \text { then } \quad t=0 ; \\
& \text { else } \\
& \quad c t 2=\tau\left(a_{i i}-a_{j j}\right) / t 2 ; \\
& \quad t=\operatorname{sign}(c t 2) /\left(\operatorname{abs}(c t 2)+\left(1+\sqrt{1+c t 2^{2}}\right) ;\right. \\
& \text { end } \\
& c s=1 / \sqrt{1+t^{2}} ; \quad s n=t / \sqrt{1+t^{2}} ; \\
& c 1=(\rho \cdot c s-\xi \cdot s n) / \tau ; \quad s 1=(\rho \cdot s n+\xi \cdot c s) / \tau ; \\
& c 2=(\rho \cdot c s+\xi \cdot s n) / \tau ; \quad s 2=(\rho \cdot s n-\xi \cdot c s) / \tau ; \\
& \delta_{i}=\left(b_{i j} / \tau-s 1\right)\left(b_{i j} / \tau+s 1\right) a_{i i}+\left(2 c 1 a_{i j}+s 2 a_{j j}\right) s 2 ; \\
& \delta_{j}=\left(s 2-b_{i j} / \tau\right)\left(s 2+b_{i j} / \tau\right) a_{j j}+\left(2 c 2 a_{i j}-s 1 a_{i i}\right) s 1 ; \\
& a_{i j}^{\prime}=(c 1 c 2-s 1 s 2) a_{i j}+\left(c 2 s 2 a_{j j}-c 1 s 1 a_{i i}\right) ; \quad a_{j i}^{\prime}=a_{i j}^{\prime} ; \\
& b_{i j}^{\prime}=0 ; \quad b_{j i}^{\prime}=b_{i j}^{\prime} ; \quad a_{i i}^{\prime}=a_{i i}+\delta_{i} ; \quad a_{i j}^{\prime}=a_{j j}-\delta_{j} ; \\
& \text { for } k=1, \ldots, n, k \neq i, j \quad \text { do } \\
& \quad a_{k i}^{\prime}=c 1 \cdot a_{k i}+s 2 \cdot a_{k j} ; \quad b_{k i}^{\prime}=c 1 \cdot b_{k i}+s 2 \cdot b_{k j} ; \quad a_{i k}^{\prime}=a_{k i}^{\prime} ; \quad b_{i k}^{\prime}=b_{k j}^{\prime} ; \\
& \quad a_{k j}^{\prime}=c 2 \cdot a_{k j}-s 1 \cdot a_{k i} ; \quad b_{k j}^{\prime}=c 2 \cdot b_{k j}-s 1 \cdot b_{k i} ; \quad a_{j k}^{\prime}=a_{k j}^{\prime} ; \quad b_{j k}^{\prime}=b_{k j}^{\prime} ; \\
& \text { endfor }
\end{aligned}
$$

endif

## New Algorithms: Based on $L L^{T}$ and $R R^{T}$ Factorizations

Consider the Cholesky foctorization of $\hat{B}: \hat{B}=\hat{L} \hat{L}^{T}$,

$$
\begin{aligned}
\hat{B} & =\hat{\hat{L}} \\
{\left[\begin{array}{cc}
1 & b_{i j} \\
b_{i j} & 1
\end{array}\right] } & =\left[\begin{array}{ll}
1 & 0 \\
a & c
\end{array}\right]\left[\begin{array}{cc}
1 & a \\
0 & c
\end{array}\right]=\left[\begin{array}{cc}
1 & a \\
a & a^{2}+c^{2}
\end{array}\right] .
\end{aligned}
$$

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Assuming $c>0$, one obtains $a=b_{i j}, \quad c=\sqrt{1-b_{i j}^{2}}$, hence

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Assuming $c>0$, one obtains $a=b_{i j}, \quad c=\sqrt{1-b_{i j}^{2}}$, hence

$$
\hat{L}=\left[\begin{array}{cc}
1 & 0 \\
b_{i j} & \sqrt{1-b_{i j}^{2}}
\end{array}\right], \quad \hat{L}^{-1}=\left[\begin{array}{cc}
1 & 0 \\
-\frac{b_{i j}}{\sqrt{1-b_{i j}^{2}}} & \frac{1}{\sqrt{1-b_{i j}^{2}}}
\end{array}\right] .
$$

## New Algorithms: Based on $L L^{T}$ and $R R^{T}$ Factorizations

Consider the Cholesky foctorization of $\hat{B}: \hat{B}=\hat{L} \hat{L}^{T}$,

$$
\begin{aligned}
\hat{B} & =\hat{\hat{L}} \\
{\left[\begin{array}{cc}
1 & b_{i j} \\
b_{i j} & 1
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1 & 0 \\
a & c
\end{array}\right]\left[\begin{array}{cc}
1 & a \\
0 & c
\end{array}\right]=\left[\begin{array}{cc}
1 & a \\
a & a^{2}+c^{2}
\end{array}\right] .
\end{aligned}
$$

Assuming $c>0$, one obtains $a=b_{i j}, \quad c=\sqrt{1-b_{i j}^{2}}$, hence

$$
\hat{L}=\left[\begin{array}{cc}
1 & 0 \\
b_{i j} & \sqrt{1-b_{i j}^{2}}
\end{array}\right], \quad \hat{L}^{-1}=\left[\begin{array}{cc}
1 & 0 \\
-\frac{b_{i j}}{\sqrt{1-b_{i j}^{2}}} & \frac{1}{\sqrt{1-b_{i j}^{2}}}
\end{array}\right] .
$$

Let the first transformation be

$$
\hat{F}_{1}=\hat{L}^{-T}, \quad \text { then } \quad \hat{F}_{1}^{T} \hat{B} \hat{F}_{1}=I_{2}
$$

## The Algorithm Based on $L L^{T}$ Factorization

$$
\hat{F}_{1}^{T} \hat{A} \hat{F}_{1}=\left[\begin{array}{cc}
a_{i i} & \frac{a_{i j}-b_{j i} a_{i j}}{\sqrt{1-b_{i j}^{2}}} \\
\frac{a_{i j}-b_{i j a i}}{\sqrt{1-b_{i j}^{i}}} & a_{j j}-\frac{2 a_{i j}-\left(a_{i} i a_{j} j b_{i j}\right.}{1-b_{i j}} b_{i j}
\end{array}\right] .
$$

## The Algorithm Based on $L L^{T}$ Factorization

$$
\hat{F}_{1}^{T} \hat{A} \hat{F}_{1}=\left[\begin{array}{cc}
a_{i i} & \frac{a_{i j}-b_{i j} a_{i i}}{\sqrt{1-b_{i j}^{2}}} \\
\frac{a_{i j}-b_{i j} a_{i i}}{\sqrt{1-b_{i j}^{2}}} & a_{j j}-\frac{2 a_{i j}-\left(a_{i i}+a_{j j}\right) b_{i j}}{1-b_{i j}^{2}} b_{i j}
\end{array}\right] .
$$

The final $\hat{F}$,

$$
\hat{F}=\hat{F}_{1} \hat{R}, \quad \hat{R} \text { is Jacobi rotation which diagonalizes } \hat{F}_{1}^{T} \hat{A} \hat{F}_{1}
$$

Its angle $\vartheta$ is determined by the formula

## The Algorithm Based on $L L^{T}$ Factorization

$$
\tan (2 \vartheta)=\frac{2\left(a_{i j}-b_{i j} a_{i i}\right) \sqrt{1-b_{i j}^{2}}}{a_{i i}-a_{j j}+2\left(a_{i j}-b_{i j} a_{i i}\right) b_{i j}}, \quad-\frac{\pi}{4} \leq \vartheta \leq \frac{\pi}{4} .
$$

## The Algorithm Based on $L L^{T}$ Factorization

$$
\tan (2 \vartheta)=\frac{2\left(a_{i j}-b_{i j} a_{i i}\right) \sqrt{1-b_{i j}^{2}}}{a_{i i}-a_{j j}+2\left(a_{i j}-b_{i j} a_{i i}\right) b_{i j}}, \quad-\frac{\pi}{4} \leq \vartheta \leq \frac{\pi}{4} .
$$

The transformation formulas for the diagonal elements of $A$ read

$$
\begin{align*}
a_{i i}^{\prime} & =a_{i i}+\tan \vartheta \cdot \frac{a_{i j}-a_{i i} b_{i j}}{\sqrt{1-b_{i j}^{2}}}  \tag{1}\\
a_{j j}^{\prime} & =a_{j j}-\frac{2 a_{i j} b_{i j}-b_{i j}^{2}\left(a_{i i}+a_{j j}\right)}{1-b_{i j}^{2}}-\tan \vartheta \cdot \frac{a_{i j}-a_{i i} b_{i j}}{\sqrt{1-b_{i j}^{2}}} \tag{2}
\end{align*}
$$

## The Algorithm Based on $L L^{T}$ Factorization

$$
\tan (2 \vartheta)=\frac{2\left(a_{i j}-b_{i j} a_{i i}\right) \sqrt{1-b_{i j}^{2}}}{a_{i i}-a_{j j}+2\left(a_{i j}-b_{i j} a_{i i}\right) b_{i j}}, \quad-\frac{\pi}{4} \leq \vartheta \leq \frac{\pi}{4} .
$$

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$$
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a_{i i}^{\prime} & =a_{i i}+\tan \vartheta \cdot \frac{a_{i j}-a_{i i} b_{i j}}{\sqrt{1-b_{i j}^{2}}}  \tag{1}\\
a_{j j}^{\prime} & =a_{j j}-\frac{2 a_{i j} b_{i j}-b_{i j}^{2}\left(a_{i i}+a_{j j}\right)}{1-b_{i j}^{2}}-\tan \vartheta \cdot \frac{a_{i j}-a_{i i} b_{i j}}{\sqrt{1-b_{i j}^{2}}} \tag{2}
\end{align*}
$$

If $a_{i j}=a_{j j}, a_{i j}=a_{i i} b_{i j}$ then $\vartheta$ is determined from $0 / 0$, so we choose $\vartheta=0$. In this case $a_{i i}^{\prime}$ and $a_{j j}^{\prime}$ reduce to $a_{i i}$ and $a_{j j}$, respectively.

## The Algorithm Based on $L L^{T}$ Factorization

This leads to a simpler matrix

$$
\begin{aligned}
\hat{Z} & =\frac{1}{\sqrt{1-b_{i j}^{2}}}\left[\begin{array}{cc}
\sqrt{1-b_{i j}^{2}} & -b_{i j} \\
0 & 1
\end{array}\right]\left[\begin{array}{cc}
c_{\vartheta} & -s_{\vartheta} \\
s_{\vartheta} & c_{\vartheta}
\end{array}\right] \\
& =\frac{1}{\sqrt{1-b_{i j}^{2}}}\left[\begin{array}{cc}
c_{\tilde{\vartheta}} & -s_{\tilde{\vartheta}} \\
s_{\vartheta} & c_{\vartheta}
\end{array}\right],
\end{aligned} \begin{aligned}
& c_{\tilde{\vartheta}}=c_{\vartheta} \sqrt{1-b_{i j}^{2}}-s_{\vartheta} b_{i j}, \\
& s_{\tilde{\vartheta}}=c_{\vartheta} b_{i j}+s_{\vartheta} \sqrt{1-b_{i j}^{2}} .
\end{aligned}
$$

It is easy to check that $c_{\tilde{\vartheta}}^{2}+s_{\tilde{\vartheta}}^{2}=1$.

## Algorithm $L L^{T} J$

```
select the pivot pair \((i, j)\)
if \(a_{i j} \neq 0\) or \(b_{i j} \neq 0\) then
    \(\beta=b_{i j}, \quad \tau=\operatorname{sqrt}((1+\beta)(1-\beta)) ; \quad \alpha=a_{i j}-\beta a_{i i} ;\)
    if \(\alpha=0 \quad\) then \(t=0\);
    else \(c t 2=\left(0.5\left(a_{i i}-a_{j j}\right)+\alpha \beta\right) /(\alpha \tau)\);
        \(t=\operatorname{sign}(c t 2) /\left(\operatorname{abs}(c t 2)+\operatorname{sqrt}\left(1+c t 2^{2}\right)\right) ;\)
    endif
    \(c s=1 / \operatorname{sqrt}\left(1+t^{2}\right) ; \quad s n=t / \operatorname{sqrt}\left(1+t^{2}\right) ;\)
    \(c 1=c s-s n \beta / \tau ; \quad s 1=s n+c s \beta / \tau ; \quad c 2=c s / \tau ; \quad s 2=s n / \tau ;\)
    \(\delta_{i}=t \alpha / \tau ; \quad \delta_{j}=\left(t \alpha+(\beta / \tau) \cdot\left(2 a_{i j}-\left(a_{i i}+a_{j j}\right) \beta\right)\right) / \tau ;\)
    \(a_{i j}^{\prime}=(c 1 c 2-s 1 s 2) a_{i j}+\left(c 2 s 2 a_{j j}-c 1 s 1 a_{i i}\right) ; \quad a_{j i}^{\prime}=a_{i j}^{\prime}\);
    \(b_{i j}^{\prime}=(c 1 c 2-s 1 s 2) \beta+(c 2 s 2-c 1 s 1) ; \quad b_{j i}^{\prime}=b_{i j}^{\prime} ;\)
    \(a_{i i}^{\prime}=a_{i i}+\delta_{i} ; \quad a_{j}^{\prime}=a_{j j}-\delta_{j}\);
    for \(k=1, \ldots, n, k \neq i, j\) do
        \(a_{k i}^{\prime}=c 1 \cdot a_{k i}+s 2 \cdot a_{k j} ; \quad b_{k i}^{\prime}=c 1 \cdot b_{k i}+s 2 \cdot b_{k j} ; \quad a_{i k}^{\prime}=a_{k i}^{\prime} ; \quad b_{i k}^{\prime}=b_{k i}^{\prime}\)
        \(a_{k j}^{\prime}=c 2 \cdot a_{k j}-s 1 \cdot a_{k i} ; \quad b_{k j}^{\prime}=c 2 \cdot b_{k j}-s 1 \cdot b_{k i} ; \quad a_{j k}^{\prime}=a_{k j}^{\prime} ; \quad b_{j k}^{\prime}=b_{k j}^{\prime} ;\)
    endfor
endif
```


## The Algorithm Based on $R R^{T}$ Factorization

$$
\begin{aligned}
\hat{B} & =\hat{R} \\
{\left[\begin{array}{cc}
1 & b_{i j} \\
b_{i j} & 1
\end{array}\right] } & =\left[\begin{array}{ll}
c & a \\
0 & 1
\end{array}\right]\left[\begin{array}{cc}
c & 0 \\
a & 1
\end{array}\right]=\left[\begin{array}{cc}
a^{2}+c^{2} & a \\
a & 1
\end{array}\right] .
\end{aligned}
$$

Assuming positive $c$, one obtains $a=b_{i j}, c=\sqrt{1-b_{i j}^{2}}$, hence

$$
\hat{R}=\left[\begin{array}{cc}
\sqrt{1-b_{i j}^{2}} & b_{i j} \\
0 & 1
\end{array}\right] \quad \text { and } \quad \hat{R}^{-1}=\left[\begin{array}{cc}
\frac{1}{\sqrt{1-b_{i j}^{2}}} & -\frac{b_{i j}}{\sqrt{1-b_{i j}^{2}}} \\
0 & 1
\end{array}\right] .
$$

## The Algorithm Based on $R R^{T}$ Factorization

$$
\begin{aligned}
\hat{B} & =\hat{R} \hat{R^{T}} \\
{\left[\begin{array}{cc}
1 & b_{i j} \\
b_{i j} & 1
\end{array}\right] } & =\left[\begin{array}{cc}
c & a \\
0 & 1
\end{array}\right]\left[\begin{array}{cc}
c & 0 \\
a & 1
\end{array}\right]=\left[\begin{array}{cc}
a^{2}+c^{2} & a \\
a & 1
\end{array}\right] .
\end{aligned}
$$

Assuming positive $c$, one obtains $a=b_{i j}, c=\sqrt{1-b_{i j}^{2}}$, hence

$$
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\frac{1}{\sqrt{1-b_{i j}^{2}}} & -\frac{b_{i j}}{\sqrt{1-b_{i j}^{2}}} \\
0 & 1
\end{array}\right] .
$$

Let $\hat{F}_{2}=\hat{R}^{-T}$, then $\hat{F}_{2}^{T} \hat{B} \hat{F}_{2}=I_{2}$

## The Algorithm Based on $R R^{T}$ Factorization

$$
\hat{F}_{2}^{T} \hat{A} \hat{F}_{2}=\left[\begin{array}{cc}
a_{i i}-\frac{2 a_{i j}-\left(a_{i j}+a_{j j}\right) b_{i j}}{1-b_{i j}^{2}} & b_{i j} \\
\frac{a_{i j}-b_{i j} a_{j j}}{\sqrt{1-b_{i j}^{2}}} \\
\sqrt{1-b_{i j}^{2}} & a_{j j}
\end{array}\right] .
$$

The final $\hat{F}$,

$$
\hat{F}=\hat{F}_{2} \hat{\jmath}, \quad \hat{\jmath} \text { is Jacobi rotation which diagonalizes } \hat{F}_{2}^{T} \hat{A} \hat{F}_{2}
$$

Its angle $\vartheta$ is determined by the formula:

$$
\tan (2 \vartheta)=\frac{2\left(a_{i j}-b_{i j} a_{j j}\right) \sqrt{1-b_{i j}^{2}}}{a_{i i}-a_{i j}-2\left(a_{i j}-b_{i j} a_{j j}\right) b_{i j}}, \quad-\frac{\pi}{4} \leq \vartheta \leq \frac{\pi}{4}
$$

## The Algorithm Based on $R R^{T}$ Factorization

The transformation formulas for the diagonal elements of $A$ :

$$
\begin{aligned}
a_{i i}^{\prime} & =a_{i i}-\frac{2 a_{i j}-\left(a_{i j}+a_{j j}\right) b_{i j}}{1-b_{i j}^{2}} b_{i j}+\tan \vartheta \cdot \frac{a_{i j}-a_{j j} b_{i j}}{\sqrt{1-b_{i j}^{2}}} \\
a_{j j}^{\prime} & =a_{j j}-\tan \vartheta \cdot \frac{a_{i j}-a_{j j} b_{i j}}{\sqrt{1-b_{i j}^{2}}}
\end{aligned}
$$

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a_{j j}^{\prime} & =a_{j j}-\tan \vartheta \cdot \frac{a_{i j}-a_{j j} b_{i j}}{\sqrt{1-b_{i j}^{2}}}
\end{aligned}
$$

If $a_{i i}=a_{j j}, a_{i j}=a_{j j} b_{i j}$ then we choose $\vartheta=0$ and then $a_{i j}^{\prime}$ and $a_{j j}^{\prime}$ reduce to $a_{i i}$ and $a_{j j}$, respectively.

## The Algorithm Based on $R R^{T}$ Factorization

This leads to the transformation matrix

$$
\begin{aligned}
\hat{Z} & =\frac{1}{\sqrt{1-b_{i j}^{2}}}\left[\begin{array}{cc}
1 & 0 \\
-b_{i j} & \sqrt{1-b_{i j}^{2}}
\end{array}\right]\left[\begin{array}{cc}
c_{\vartheta} & -s_{\vartheta} \\
s_{\vartheta} & c_{\vartheta}
\end{array}\right] \\
& =\frac{1}{\sqrt{1-b_{i j}^{2}}}\left[\begin{array}{cc}
c_{\vartheta} & -s_{\vartheta} \\
s_{\tilde{\vartheta}} & c_{\tilde{\vartheta}}
\end{array}\right],
\end{aligned} \begin{aligned}
& c_{\tilde{\vartheta}}=c_{\vartheta} \sqrt{1-b_{i j}^{2}}+s_{\vartheta} b_{i j}, \\
& s_{\tilde{\vartheta}}=s_{\vartheta} \sqrt{1-b_{i j}^{2}}-c_{\vartheta} b_{i j} .
\end{aligned}
$$

It is easy to check that $c_{\tilde{\vartheta}}^{2}+s_{\tilde{\vartheta}}^{2}=1$.

## Algorithm $R R^{T} J$

select the pivot pair $(i, j)$
if $a_{i j} \neq 0$ or $b_{i j} \neq 0$ then

$$
\begin{aligned}
& \beta=b_{i j}, \tau=\operatorname{sqrt}((1+\beta)(1-\beta)) ; \quad \alpha=a_{i j}-\beta a_{j j} ; \\
& \text { if } \alpha=0 \quad \text { then } \quad t=0 ; \\
& \text { else } \quad c t 2=\left(0.5\left(a_{i i}-a_{j j}\right)-\alpha \beta\right) /(\alpha \tau) ; \\
& \quad t=\operatorname{sign}(c t 2) /\left(\operatorname{abs}(c t 2)+\operatorname{sqrt}\left(1+c t 2^{2}\right)\right) ;
\end{aligned}
$$

endif

$$
\begin{aligned}
& c s=1 / \text { sqrt }\left(1+t^{2}\right) ; \quad \text { sn }=t / \text { sqrt }\left(1+t^{2}\right) ; \\
& c 1=c s / \tau ; \quad s 1=s n / \tau ; \quad c 2=c s+s n \beta / \tau ; \quad s 2=s n-c s \beta / \tau ; \\
& \delta_{j}=t \alpha / \tau ; \quad \delta_{i}=\left(t \alpha-(\beta / \tau) \cdot\left(2 a_{i j}-\left(a_{i i}+a_{j j}\right) \beta\right)\right) / \tau ; \\
& a_{i j}^{\prime}=(c 1 c 2-s 1 s 2) a_{i j}+\left(c 2 s 2 a_{j j}-c 1 s 1 a_{i j}\right) ; a_{j i}^{\prime}=a_{i j}^{\prime} ; \\
& b_{i j}^{\prime}=(c 1 c 2-s 1 s 2) \beta+(c 2 s 2-c 1 s 1) ; \quad b_{j i}^{\prime}=b_{i j}^{\prime} ; \\
& a_{i i j}^{\prime}=a_{i i}+\delta_{i j} ; \quad a_{j}^{\prime}=a_{j j}-\delta_{j} ; \\
& \text { for } k=1, \ldots, n, k \neq i, j \quad \text { do } \\
& \quad a_{k i}^{\prime}=c 1 \cdot a_{k i}+s 2 \cdot a_{k j} ; \quad b_{k i}^{\prime}=c 1 \cdot b_{k i}+s 2 \cdot b_{k j} ; \quad a_{i k}^{\prime}=a_{k k}^{\prime} ; \quad b_{i k}^{\prime}=b_{k i}^{\prime} \\
& a_{k j}^{\prime}=c 2 \cdot a_{k j}-s 1 \cdot a_{k i} ; \quad b_{k j}^{\prime}=c 2 \cdot b_{k j}-s 1 \cdot b_{k i} ; \quad a_{j k}^{\prime}=a_{k j}^{\prime} ; \quad b_{j k}^{\prime}=b_{k j}^{\prime} ;
\end{aligned}
$$

endfor
endif

## Definition of a Hybrid and a General Method

## Definition

Let $\mathcal{H}$ denote a collection of Jacobi methods for PGEP $A x=\lambda B x$ which satisfy the following two rules:
(1) at step $k, \hat{A}^{(k)}$ is diagonalized and $\hat{B}^{(k)}$ is transformed to $I_{2}$,
(2) at least one diagonal element of $\hat{F}_{k}$ is not smaller than $\sqrt{2} / 2$.

An element of $\mathcal{H}$ is called a general PGEP Jacobi method.
A hybrid Jacobi method is any method from $\mathcal{H}$ that uses at each step either the $\mathrm{HZ}, L L^{T} J$ or $R R^{T} J$ algorithm.

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An element of $\mathcal{H}$ is called a general PGEP Jacobi method.
A hybrid Jacobi method is any method from $\mathcal{H}$ that uses at each step either the $\mathrm{HZ}, L L^{T} J$ or $R R^{T} J$ algorithm.

In this definition the pivot strategy is not specified, hence any can be used. If a method uses only the $\mathrm{HZ}\left(L L^{T} J, R R^{T} J\right)$ algorithm, it will be called the $\mathrm{HZ}\left(L L^{T} J, R R^{T} J\right)$ method.

## Some Remarks

- It is easy to show that $\mathrm{HZ}, L L^{T} J$ and $R R^{T} J$ methods belong to $\mathcal{H}$


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- Algorithms based on $L L^{T}$ and $R R^{T}$ factorizations are called $L L^{T} J$ and $R R^{T} J$ algorithm, because $L L^{T}$ and $R R^{T}$ factorizations are followed by one step of the standard Jacobi method


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- It is easy to show that $\mathrm{HZ}, L L^{T} J$ and $R R^{T} J$ methods belong to $\mathcal{H}$
- Algorithms based on $L L^{T}$ and $R R^{T}$ factorizations are called $L L^{T} J$ and $R R^{T} J$ algorithm, because $L L^{T}$ and $R R^{T}$ factorizations are followed by one step of the standard Jacobi method
- The general (PGEP) Jacobi method can use at each step any conceivable algorithm which satisfies the above two rules. For example, it can use the FL method combined with normalization of the elements of $B$


## Some Remarks

- All real algorithms have the form

$$
\hat{Z}=\frac{1}{\sqrt{1-b_{i j}^{2}}}\left[\begin{array}{cc}
\cos \phi & -\sin \phi \\
\cos \psi & \sin \psi
\end{array}\right]
$$

This follows from a result of Gose (ZAMM 59, 1979), who found the general form of a matrix $\hat{Z}$ which diagonalizes a $\hat{B} \succ O$ via the congruence transformation $\hat{B} \mapsto \hat{Z}^{\top} \hat{B} \hat{Z}$.

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If we assume $\left.b_{11}=\cdots=b_{n n}\right) 1$ and the same for $\hat{Z}^{T} \hat{B} \hat{Z}$, then this form of $\hat{Z}$ is just the Gose's theorem.

## Thank you for your attention

