

Multipartite Secret Sharing

Laszlo Csirmaz

UTIA, Prague
Rényi Institute, Budapest

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Contents

- 1 Secret sharing
- 2 Capped structures
- 3 Bipartite and tripartite ideal structures
- 4 Some proofs
- 5 Acknowledgments

Secret sharing by groups

- Participants are in disjoint groups

$$P = P_1 \cup P_2 \cup \dots \cup P_m.$$

Sometimes we call them *departments*.

- Members of each group play the same role
any participant can be replaced by any other member from the same group.
- Interesting only if there are few groups and several members in each group.
- Many unsolved problems
even for the bipartite (two groups) case.

Definitions

- **Access structure**
is the collection of qualified sets.
- **Complexity**
is the maximal relative share size; it is at least 1
- **Ideal structures**
are the ones with minimal complexity 1.
- **κ -ideal structures**
are where the entropy method gives the lower bound 1 on the complexity (not necessarily ideal).

Theorem (Brickell & Davenport – informal)

κ -ideal access structures and matroids are in a one-to-one correspondence.



The “cap” theorem

Theorem (Csirmaz & Matúš & Padró – informal)

Multipartite κ -ideal structures are the same as “capped” structures.

- 1 For $m = 1$ “capped” structures are just the threshold ones.
- 2 Recipe to list / generate / recognize all such structures.
- 3 For $m = 1$, $m = 2$, and $m = 3$ “capped” structures are linearly representable.

Corollary

We have a complete description of all ideal tripartite access structures.

- 4 For $m = 4$ there is a κ -ideal structure which is not ideal.

Contents

- 1 Secret sharing
- 2 Capped structures**
- 3 Bipartite and tripartite ideal structures
- 4 Some proofs
- 5 Acknowledgments

Capped structures



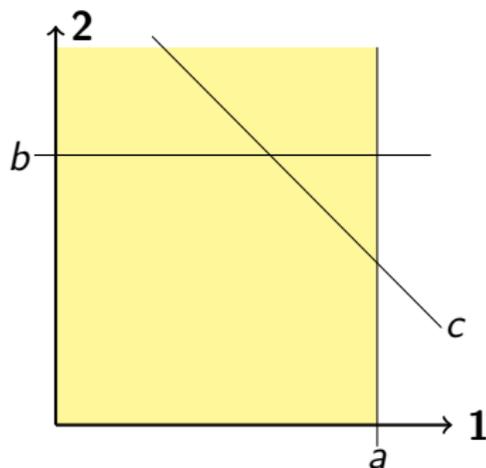
Capped structures

Each subset A of the groups (departments) has a **cap** $f(A)$.

Mnemonic: the power of the coalition A of some departments is limited to $f(A)$ counts.

Example:

Departments: $\{1, 2\}$; $f(1) = a$, $f(2) = b$, $f(12) = c$:



cap a for department **1**

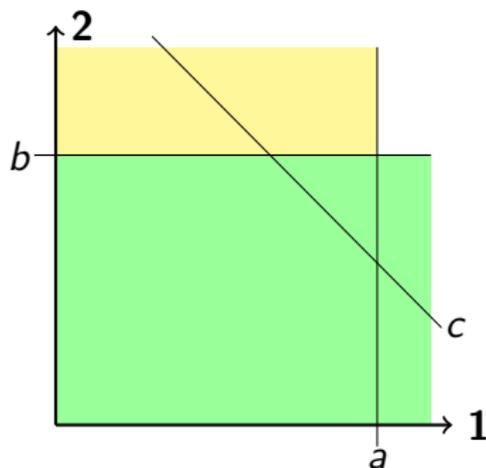
Capped structures

Each subset A of the groups (departments) has a **cap** $f(A)$.

Mnemonic: the power of the coalition A of some departments is limited to $f(A)$ counts.

Example:

Departments: $\{1, 2\}$; $f(1) = a$, $f(2) = b$, $f(12) = c$:



cap b for department 2

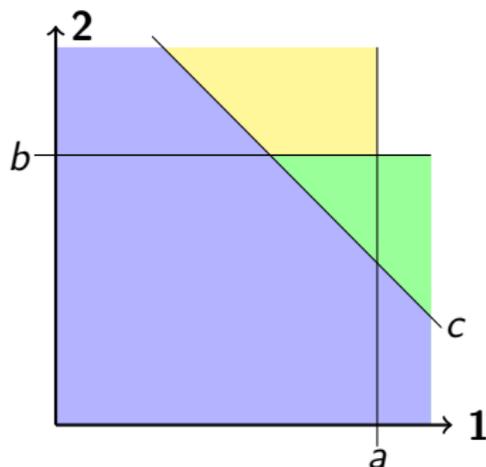
Capped structures

Each subset A of the groups (departments) has a **cap** $f(A)$.

Mnemonic: the power of the coalition A of some departments is limited to $f(A)$ counts.

Example:

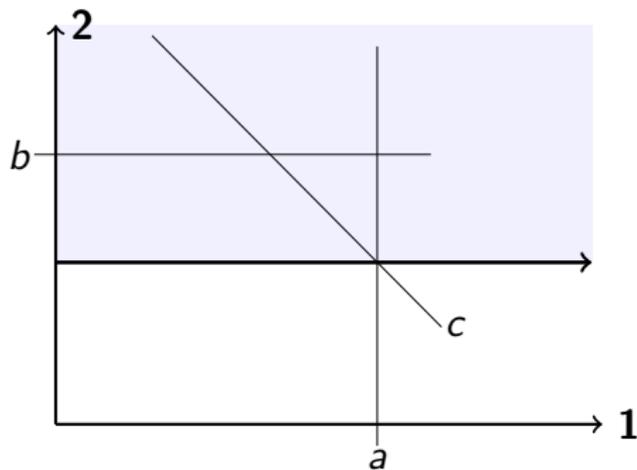
Departments: $\{1, 2\}$; $f(1) = a$, $f(2) = b$, $f(12) = c$:



cap c for both departments

Hitting the cap c

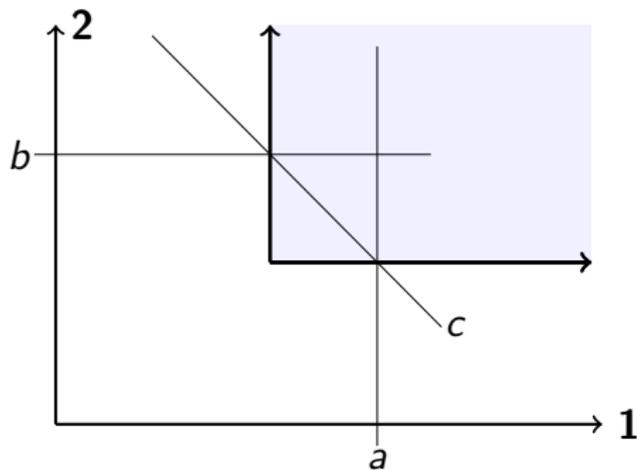
As $f(\mathbf{1}) = a$, there must be at least $c - a$ members from group **2**.



Hitting the cap c

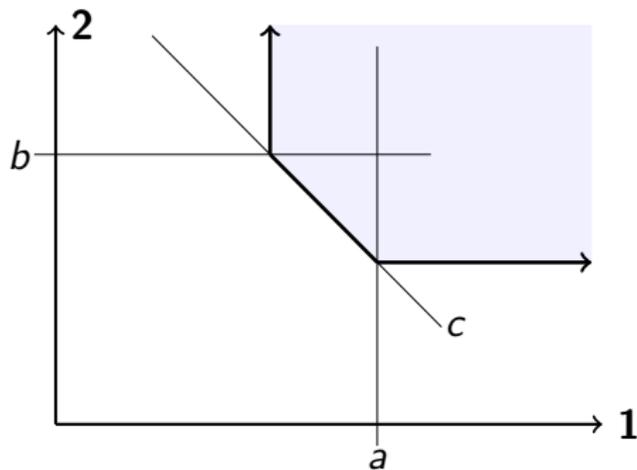
As $f(1) = a$, there must be at least $c - a$ members from group **2**.

As $f(2) = b$, there must be at least $c - b$ members from group **1**.



Hitting the cap c

As $f(1) = a$, there must be at least $c - a$ members from group **2**.
As $f(2) = b$, there must be at least $c - b$ members from group **1**.
And at least c members from the two groups together.





The cap function f

Participants are in m disjoint groups (departments)

$$P = P_1 \cup P_2 \cup \dots \cup P_m.$$

For each subset A of the groups $f(A)$ is the “cap” of A so that

- 1 $f(\emptyset) = 0$, otherwise $f(A)$ is a positive integer,
- 2 f is monotonic: $f(A) \leq f(A \cup B)$,
- 3 f is submodular:

$$f(A) + f(B) \geq f(A \cap B) + f(A \cup B).$$

Otherwise there is no way to hit the the cap $f(A \cup B)$.



Capped structures

In secret sharing a capped access structure is defined by

- the set of participants P who are in m disjoint groups:

$$P = P_1 \cup P_2 \cup \dots \cup P_m,$$

- the cap function $f(A)$ defined for each subset of the groups,
- an upward closed collection of group subsets:

$$\mathcal{A} = \{A_1, A_2, \dots, A_t\}$$

(if $B \supset A_i$, then B is also in \mathcal{A}).

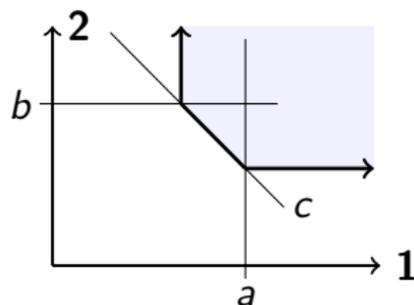
Definition ( Capped access structure)

A subset of participants is qualified if and only if they hit the cap $f(A_j)$ for some $A_j \in \mathcal{A}$.

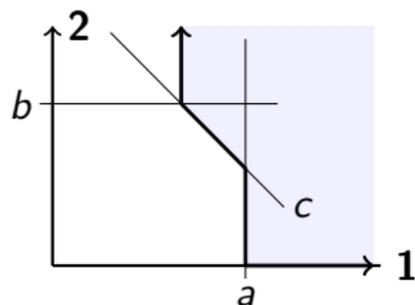
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- 2 Capped structures
- 3 Bipartite and tripartite ideal structures**
- 4 Some proofs
- 5 Acknowledgments

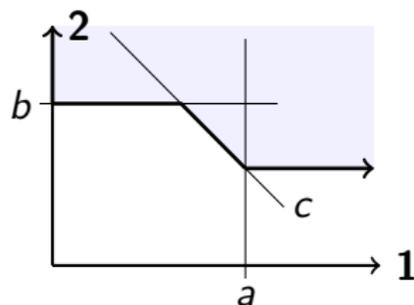
Case of two departments 1 and 2



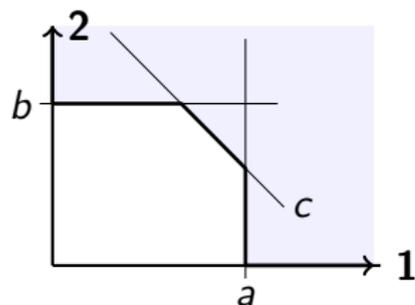
$$\mathcal{A} = \{12\}$$



$$\mathcal{A} = \{1, 12\}$$



$$\mathcal{A} = \{2, 12\}$$



$$\mathcal{A} = \{1, 2, 12\}$$

Case of three departments 1, 2, 3

Seven cap values:

$$f(\mathbf{123})$$

$$f(\mathbf{12}) \quad f(\mathbf{13}) \quad f(\mathbf{23})$$

$$f(\mathbf{1}) \quad f(\mathbf{2}) \quad f(\mathbf{3})$$

numerous possibilities for \mathcal{A} , e.g.,

$$\mathcal{A} = \{\mathbf{1}, \mathbf{12}, \mathbf{13}, \mathbf{123}\},$$

each yielding an ideal structure.



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The C-M-P theorem, main points

- Σ is a κ -ideal multipartite structure with partition π .
- The matroid M corresponds to Σ (Brickell-Davenport thm).
- Factor M by the partition to get $N = M/\pi$, an integer polymatroid on the partition groups.
Note: the ranks of N define the  values!
- M can be recovered from N uniquely (due to the multipartite symmetry).
- The secret defines a one-point extension of M (and of N) and it has rank 1. Qualified subsets are those whose rank is not increased by this extension.
- Such a one-point extension is characterized by a *modular cut* in the factor polymatroid N : this is the collection of all flats whose ranks do not increase – the collection \mathcal{A} in the examples.

Tripartite κ -ideal structures are linear

- In the tripartite case the factor polymatroid N is integer and it is on three points. Such polymatroids are known to be linear.
- **If** the one-point extension of N (by the secret) is linear, **then** M is linear. There are arbitrary large vector space representations and one can choose many “generic” elements.
- An integer polymatroid on a, b, c, d is linearly representable if and only if it satisfies all instances of the Ingleton inequality

$$0 \leq \text{ING}(a, b, c, d) = f(ab) + f(ac) + f(ad) + f(bc) + f(bd) - f(a) - f(b) - f(abc) - f(abd) - f(cd).$$
- In any polymatroid, $2 \cdot \text{ING}(a, b, c, d) + f(s) \geq 0$ where s is any of a, b, c, d .
- The one-point extension $N \cup \{s\}$ is integer with $f(s) = 1$. Thus $\text{ING}(a, b, c, d)$ is integer and at least $-1/2$, thus non-negative.

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